

## 12. Primal–dual proximal methods

- primal–dual optimality conditions
- primal–dual hybrid gradient algorithm
- monotone operators
- proximal point algorithm

# Primal and dual problem

primal:      minimize     $f(x) + g(Ax)$

dual:        maximize     $-g^*(z) - f^*(-A^T z)$

- $f$  and  $g$  are closed convex functions
- dual problem is Lagrange dual of reformulated problem

minimize     $f(x) + g(y)$   
subject to    $Ax = y$

## Optimality (Karush–Kuhn–Tucker) conditions (see pp. 5.21–5.24)

- primal feasibility:  $x \in \text{dom } f$  and  $Ax = y \in \text{dom } g$
- $(x, y)$  is a minimizer of the Lagrangian  $f(x) + g(y) + z^T (Ax - y)$ :

$$-A^T z \in \partial f(x), \quad z \in \partial g(y) \quad (\text{equivalently, } y \in \partial g^*(z))$$

# Primal–dual optimality conditions

- the optimality conditions can be written symmetrically as

$$0 \in \begin{bmatrix} 0 & A^T \\ -A & 0 \end{bmatrix} \begin{bmatrix} x \\ z \end{bmatrix} + \begin{bmatrix} \partial f(x) \\ \partial g^*(z) \end{bmatrix}$$

- second term on right-hand side denotes the product set

$$\partial f(x) \times \partial g^*(z) = \{(u, v) \mid u \in \partial f(x), v \in \partial g^*(z)\}$$

- solutions are saddle points of convex–concave function

$$f(x) - g^*(z) + z^T Ax$$

in this lecture we assume that the optimality conditions are solvable  
(a sufficient condition is that primal is solvable and Slater's condition holds)

# Outline

- primal–dual optimality conditions
- **primal–dual hybrid gradient algorithm**
- monotone operators
- proximal point algorithm

# Primal–dual hybrid gradient (PDHG) method

$$0 \in \begin{bmatrix} 0 & A^T \\ -A & 0 \end{bmatrix} \begin{bmatrix} x \\ z \end{bmatrix} + \begin{bmatrix} \partial f(x) \\ \partial g^*(z) \end{bmatrix}$$

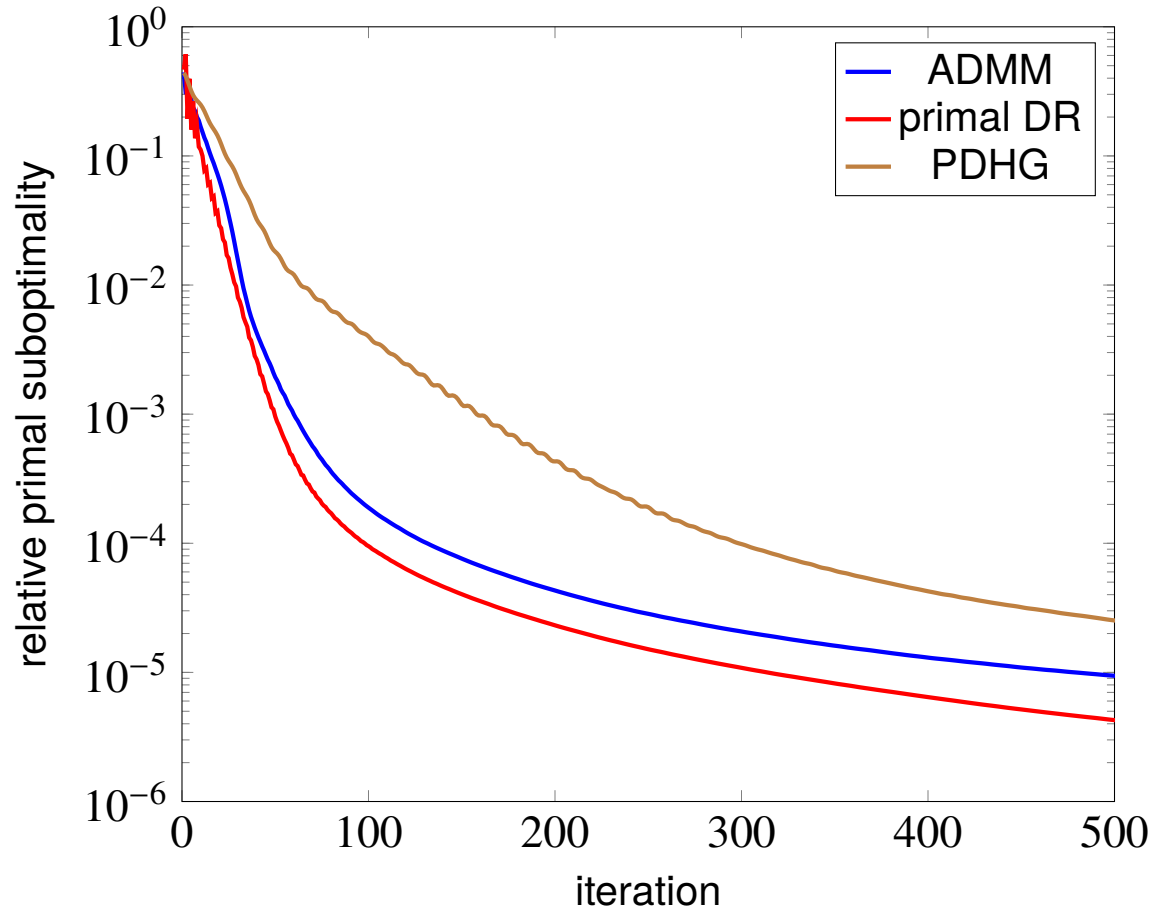
## Algorithm

$$\begin{aligned} x_{k+1} &= \text{prox}_{\tau f}(x_k - \tau A^T z_k) \\ z_{k+1} &= \text{prox}_{\sigma g^*}(z_k + \sigma A(2x_{k+1} - x_k)) \end{aligned}$$

- each iteration requires evaluations of proximal mappings of  $f$  and  $g^*$
- requires multiplications with  $A$  and  $A^T$ , but no solutions of linear equations
- primal and dual step sizes  $\tau$ ,  $\sigma$  are positive and must satisfy  $\sigma\tau\|A\|_2^2 \leq 1$

# Example

same problem as on pp. 11.20–11.24



- multiplications with  $A$  and  $A^T$  require 2-D FFTs
- with periodic boundary conditions, cost/iteration is similar for the three methods

# Douglas–Rachford method derived from PDHG

$$\text{minimize } f(x) + g(x)$$

- a special case of the standard problem on page 12.2 with  $A = I$
- apply PDHG with  $\sigma = \tau = 1$ :

$$x_{k+1} = \text{prox}_f(x_k - z_k)$$

$$z_{k+1} = \text{prox}_{g^*}(z_k + 2x_{k+1} - x_k)$$

- this is the primal–dual form of the Douglas–Rachford method on page 11.8

# PDHG derived from Douglas–Rachford method

apply the Douglas–Rachford splitting method to a reformulation of the problem:

$$\begin{array}{ll} \text{minimize} & f(x) + g(Ax) \\ & \longrightarrow \\ \text{minimize} & f(x) + g(Ax + By) \\ \text{subject to} & y = 0 \end{array}$$

- $B$  is chosen to satisfy

$$AA^T + BB^T = (1/\alpha)I \quad \text{where } 1/\alpha \geq \|A\|_2^2$$

for example,  $B = ((1/\alpha)I - AA^T)^{1/2}$

- reformulated problem is equivalent to minimizing  $\tilde{f}(x, y) + \tilde{g}(x, y)$  with

$$\tilde{f}(x, y) = f(x) + \delta_{\{0\}}(y), \quad \tilde{g}(x, y) = g(Ax + By)$$

- after simplifications, DR applied to reformulated problem will reduce to PDHG



## Proximal operators for reformulated problem

- proximal operator of  $\tilde{f}(x, y) = f(x) + \delta_{\{0\}}(y)$ :

$$\text{prox}_{\tau\tilde{f}}(x, y) = \begin{bmatrix} \text{prox}_{\tau f}(x) \\ 0 \end{bmatrix}$$

- proximal operator of  $\tilde{g}(x, y) = g(Ax + By)$  follows from page 6.8 and page 6.7:

$$\begin{aligned} \text{prox}_{\tau\tilde{g}}(x, y) &= \begin{bmatrix} x \\ y \end{bmatrix} - \alpha \begin{bmatrix} A^T \\ B^T \end{bmatrix} (Ax + By - \text{prox}_{(\tau/\alpha)g}(Ax + By)) \\ &= \begin{bmatrix} x \\ y \end{bmatrix} - \tau \begin{bmatrix} A^T \\ B^T \end{bmatrix} \text{prox}_{\sigma g^*}(\sigma(Ax + By)) \end{aligned}$$

where  $\sigma = \alpha/\tau$

- proximal operator of  $\tilde{g}^*$  follows from Moreau identity (page 6.6)

$$\text{prox}_{(\tau\tilde{g})^*}(x, y) = \tau \begin{bmatrix} A^T \\ B^T \end{bmatrix} \text{prox}_{\sigma g^*}(\sigma(Ax + By))$$

# Douglas–Rachford algorithm applied to reformulated problem

$$\text{minimize } \underbrace{f(x) + \delta_{\{0\}}(y)}_{\tilde{f}(x,y)} + \underbrace{g(Ax + By)}_{\tilde{g}(x,y)}$$

- primal–dual form of Douglas–Rachford algorithm (page 11.8) with step size  $\tau$

$$\begin{bmatrix} x_{k+1} \\ y_{k+1} \end{bmatrix} = \text{prox}_{\tau \tilde{f}} \left( \begin{bmatrix} x_k - p_k \\ y_k - q_k \end{bmatrix} \right)$$

$$\begin{bmatrix} p_{k+1} \\ q_{k+1} \end{bmatrix} = \text{prox}_{(\tau \tilde{g})^*} \left( \begin{bmatrix} p_k + 2x_{k+1} - x_k \\ q_k + 2y_{k+1} - y_k \end{bmatrix} \right)$$

- substitute expressions for proximal operators (with  $\sigma$  defined as  $\sigma = \alpha/\tau$ )

$$\begin{bmatrix} x_{k+1} \\ y_{k+1} \end{bmatrix} = \begin{bmatrix} \text{prox}_{\tau f}(x_k - p_k) \\ 0 \end{bmatrix}$$

$$\begin{bmatrix} p_{k+1} \\ q_{k+1} \end{bmatrix} = \tau \begin{bmatrix} A^T \\ B^T \end{bmatrix} \text{prox}_{\sigma g^*} \left( \sigma \begin{bmatrix} A & B \end{bmatrix} \begin{bmatrix} p_k + 2x_{k+1} - x_k \\ q_k + 2y_{k+1} - y_k \end{bmatrix} \right)$$

## First simplification

$$\begin{bmatrix} x_{k+1} \\ y_{k+1} \end{bmatrix} = \begin{bmatrix} \text{prox}_{\tau f}(x_k - p_k) \\ 0 \end{bmatrix}$$

$$\begin{bmatrix} p_{k+1} \\ q_{k+1} \end{bmatrix} = \tau \begin{bmatrix} A^T \\ B^T \end{bmatrix} \text{prox}_{\sigma g^*} \left( \sigma \begin{bmatrix} A & B \end{bmatrix} \begin{bmatrix} p_k + 2x_{k+1} - x_k \\ q_k + 2y_{k+1} - y_k \end{bmatrix} \right)$$

- from first step,  $y_k = 0$  for all  $k$  if we start with  $y_0 = 0$
- we remove the zero variable  $y_k$ :

$$x_{k+1} = \text{prox}_{\tau f}(x_k - p_k)$$

$$\begin{bmatrix} p_{k+1} \\ q_{k+1} \end{bmatrix} = \tau \begin{bmatrix} A^T \\ B^T \end{bmatrix} \text{prox}_{\sigma g^*} \left( \sigma \begin{bmatrix} A & B \end{bmatrix} \begin{bmatrix} p_k \\ q_k \end{bmatrix} + \sigma A(2x_{k+1} - x_k) \right)$$

## Second simplification

$$x_{k+1} = \text{prox}_{\tau f}(x_k - p_k)$$

$$\begin{bmatrix} p_{k+1} \\ q_{k+1} \end{bmatrix} = \tau \begin{bmatrix} A^T \\ B^T \end{bmatrix} \text{prox}_{\sigma g^*}(\sigma(Ap_k + Bq_k) + \sigma A(2x_{k+1} - x_k))$$

- from step 2:  $\begin{bmatrix} p_k \\ q_k \end{bmatrix} \in \text{range} \begin{bmatrix} A^T \\ B^T \end{bmatrix}$  for all  $k$ , if this holds for  $(p_0, q_0)$
- since  $AA^T + BB^T = (1/\alpha)I$  and  $\sigma = \alpha/\tau$ ,

$$\begin{bmatrix} p_k \\ q_k \end{bmatrix} = \tau \begin{bmatrix} A^T \\ B^T \end{bmatrix} z_k \quad \text{for a unique } z_k \text{ given by } z_k = \sigma(Ap_k + Bq_k)$$

- a change of variables  $z_k = \sigma(Ap_k + Bq_k)$  gives

$$x_{k+1} = \text{prox}_{\tau f}(x_k - \tau A^T z_k), \quad z_{k+1} = \text{prox}_{\sigma g^*}(z_k + \sigma A(2x_{k+1} - x_k))$$

this is the PDHG algorithm with  $\sigma\tau = \alpha \leq 1/\|A\|_2^2$

## PDHG with overrelaxation

$$\begin{aligned}\bar{x}_{k+1} &= \operatorname{prox}_{\tau f}(x_k - \tau A^T z_k) \\ \bar{z}_{k+1} &= \operatorname{prox}_{\sigma g^*}(z_k + \sigma A(2\bar{x}_{k+1} - x_k)) \\ \begin{bmatrix} x_{k+1} \\ z_{k+1} \end{bmatrix} &= \begin{bmatrix} x_k \\ z_k \end{bmatrix} + \rho_k \begin{bmatrix} \bar{x}_{k+1} - x_k \\ \bar{z}_{k+1} - z_k \end{bmatrix}\end{aligned}$$

- $\rho_k \in (0, 2)$
- convergence follows from convergence of Douglas–Rachford splitting method
- other types of acceleration exist for problems with strongly convex  $f$  or  $g^*$

# Outline

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- primal–dual hybrid gradient algorithm
- **monotone operators**
- proximal point algorithm

# Multivalued (set-valued) operator

**Definition:** operator  $F$  maps vectors  $x \in \mathbf{R}^n$  to sets  $F(x) \subseteq \mathbf{R}^n$

- the domain and graph of  $F$  are defined as

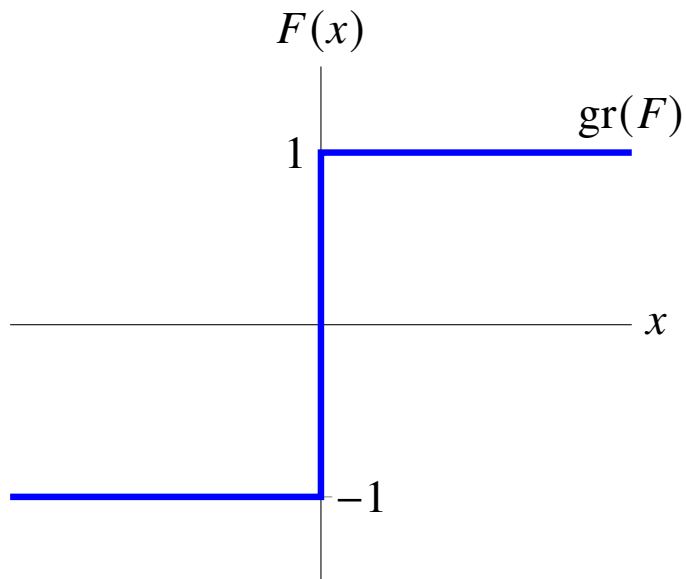
$$\text{dom } F = \{x \in \mathbf{R}^n \mid F(x) \neq \emptyset\}$$

$$\text{gr}(F) = \{(x, y) \in \mathbf{R}^n \times \mathbf{R}^n \mid x \in \text{dom } F, y \in F(x)\}$$

- if  $F(x)$  is a singleton, we write  $F(x) = y$  instead of  $F(x) = \{y\}$

**Example:** sign operator

$$F(x) = \begin{cases} -1 & x < 0 \\ [-1, 1] & x = 0 \\ 1 & x > 0 \end{cases}$$



# Transformations as operations on graph

**Inverse:**  $F^{-1}(x) = \{y \mid x \in F(y)\}$

$$\text{gr}(F^{-1}) = \begin{bmatrix} 0 & I \\ I & 0 \end{bmatrix} \text{gr}(F)$$

**Composition with scaling:**  $(\lambda F)(x) = \lambda F(x)$  and  $(F\lambda)(x) = F(\lambda x)$

$$\text{gr}(\lambda F) = \begin{bmatrix} I & 0 \\ 0 & \lambda I \end{bmatrix} \text{gr}(F), \quad \text{gr}(F\lambda) = \begin{bmatrix} (1/\lambda)I & 0 \\ 0 & I \end{bmatrix} \text{gr}(F)$$

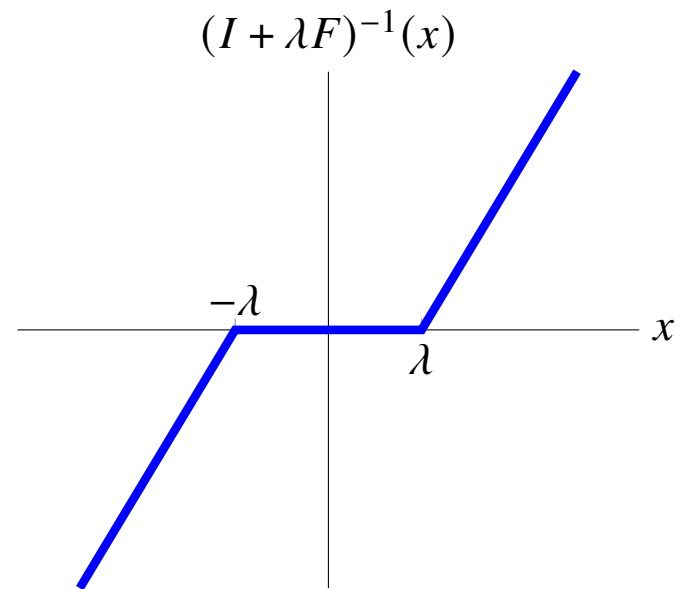
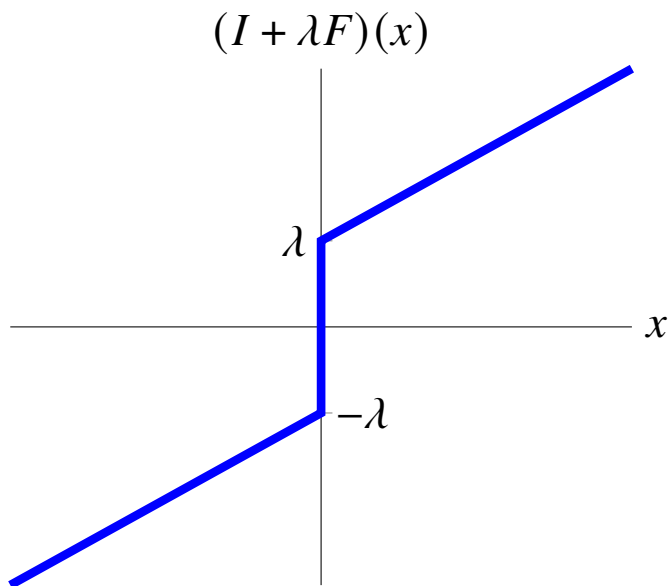
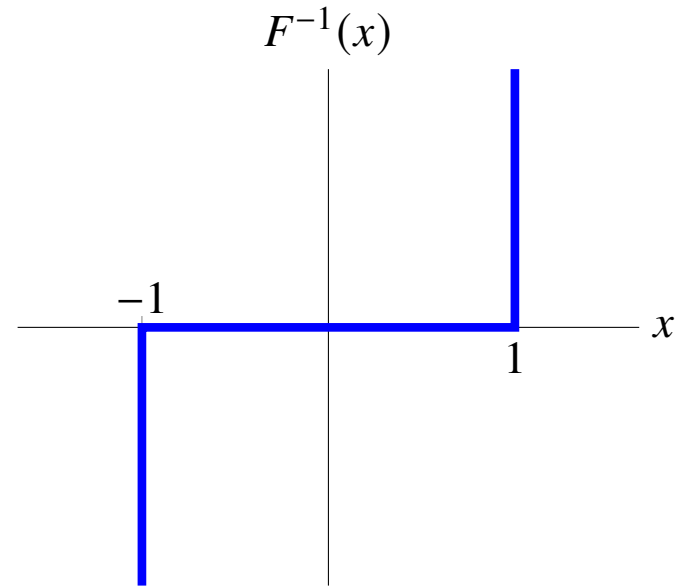
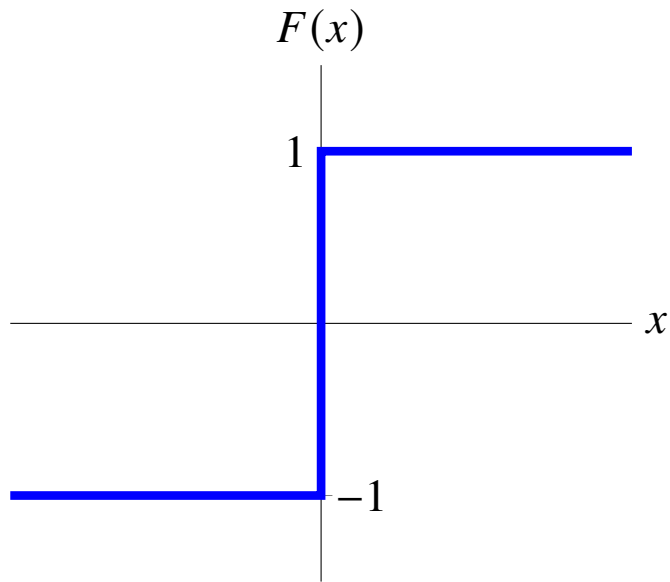
**Addition to identity:**  $(I + \lambda F)(x) = \{x + \lambda y \mid y \in F(x)\}$

$$\text{gr}(I + \lambda F) = \begin{bmatrix} I & 0 \\ I & \lambda I \end{bmatrix} \text{gr}(F)$$

note that these are all *linear* operations on the graph



# Example



# Monotone operator

**Definition:**  $F$  is a monotone operator if

$$(y - \hat{y})^T (x - \hat{x}) \geq 0 \quad \text{for all } x, \hat{x} \in \text{dom } F, y \in F(x), \hat{y} \in F(\hat{x})$$

in terms of the graph,

$$\begin{bmatrix} x - \hat{x} \\ y - \hat{y} \end{bmatrix}^T \begin{bmatrix} 0 & I \\ I & 0 \end{bmatrix} \begin{bmatrix} x - \hat{x} \\ y - \hat{y} \end{bmatrix} \geq 0 \quad \text{for all } (x, y), (\hat{x}, \hat{y}) \in \text{gr}(F)$$

**Monotone inclusion problem:** find  $x \in F^{-1}(0)$ , *i.e.*, solve

$$0 \in F(x)$$

this covers many equilibrium/optimalty conditions as special cases

# Examples

we will encounter the following three types of monotone operators

- subdifferentials  $\partial f(x)$  of convex functions  $f$
- affine monotone operators:  $F(x) = Cx + d$  is monotone if

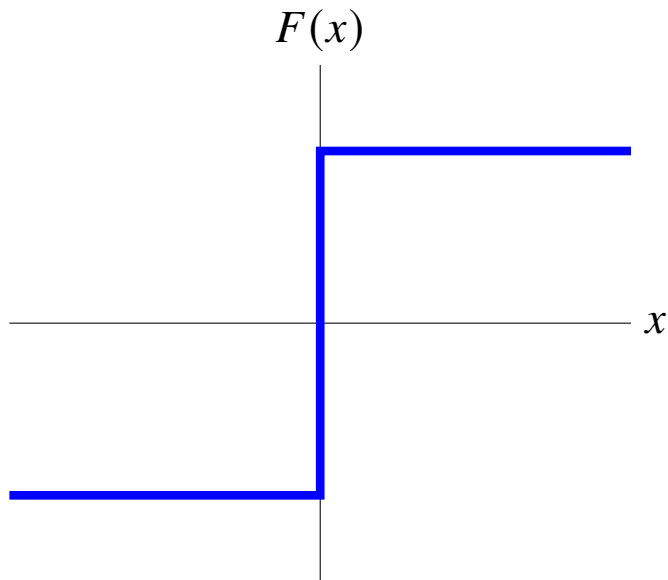
$$C + C^T \geq 0$$

- sums of the above; in particular,

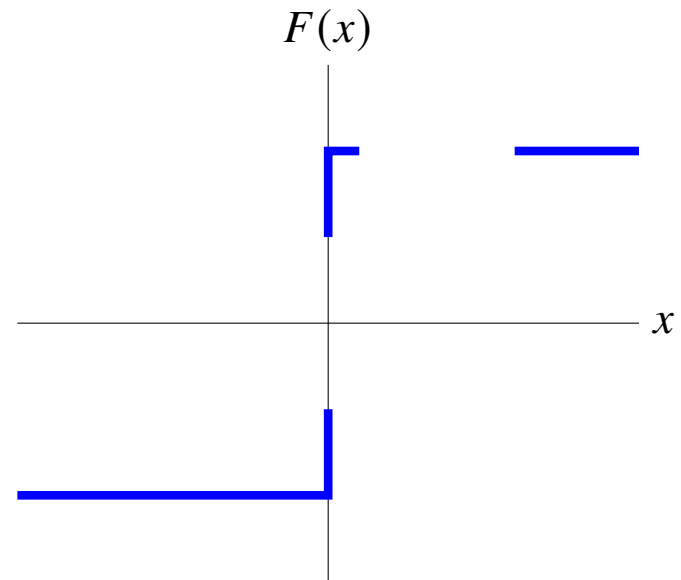
$$F(x, z) = \begin{bmatrix} 0 & A^T \\ -A & 0 \end{bmatrix} \begin{bmatrix} x \\ z \end{bmatrix} + \begin{bmatrix} \partial f(x) \\ \partial g^*(z) \end{bmatrix}$$

# Maximal monotone operator

graph is not properly contained in the graph of another monotone operator



maximal monotone



monotone, but not maximal monotone

# Conditions for maximal monotonicity

- the subdifferential of a closed convex function is maximal monotone
- affine monotone operators are maximal monotone
- (Minty's theorem) a monotone operator  $F$  is maximal monotone if and only if

$$\text{im}(I + F) := \bigcup_{x \in \text{dom } F} (x + F(x)) = \mathbf{R}^n$$

*i.e.*, for every  $y \in \mathbf{R}^n$ , there exists an  $x$  such that  $y \in x + F(x)$

- sums  $F + G$  of maximal monotone operators are not necessarily maximal  
(sufficient condition:  $\text{int dom } F \cap \text{dom } G \neq \emptyset$ )

## Coercivity (strong monotonicity)

$F$  is **coercive** with parameter  $\mu > 0$  if

$$(y - \hat{y})^T (x - \hat{x}) \geq \mu \|x - \hat{x}\|_2^2 \quad \text{for all } x, \hat{x} \in \text{dom } F, y \in F(x), \hat{y} \in F(\hat{x})$$

- $F - \mu I$  is a monotone operator
- equivalently,

$$\begin{bmatrix} x - \hat{x} \\ y - \hat{y} \end{bmatrix}^T \begin{bmatrix} -2\mu I & I \\ I & 0 \end{bmatrix} \begin{bmatrix} x - \hat{x} \\ y - \hat{y} \end{bmatrix} \geq 0 \quad \text{for all } (x, y), (\hat{x}, \hat{y}) \in \text{gr}(F)$$

### Examples

- subdifferential of strongly convex function
- affine operator  $F(x) = Ax + b$  if  $A + A^T \succ 0$  (with  $\mu = \lambda_{\min}(A + A^T)/2$ )

# Co-coercivity

$F$  is **co-coercive** with parameter  $\gamma > 0$  if  $F^{-1}$  is coercive

- $F(x)$  is single-valued for  $x \in \text{dom } F$  and

$$(F(x) - F(\hat{x}))^T (x - \hat{x}) \geq \gamma \|F(x) - F(\hat{x})\|_2^2 \quad \text{for all } x, \hat{x} \in \text{dom } F$$

- equivalently,

$$\begin{bmatrix} x - \hat{x} \\ y - \hat{y} \end{bmatrix}^T \begin{bmatrix} 0 & I \\ I & -2\gamma I \end{bmatrix} \begin{bmatrix} x - \hat{x} \\ y - \hat{y} \end{bmatrix} \geq 0 \quad \text{for all } (x, y), (\hat{x}, \hat{y}) \in \text{gr}(F)$$

- $F$  is **firmly nonexpansive** if it is co-coercive with  $\gamma = 1$

**Example:** affine operator  $F(x) = Ax + b$  with

$$A + A^T \geq 2\gamma A^T A \quad \iff \quad \|2\gamma A - I\|_2 \leq 1$$

for symmetric positive definite  $A$ , this means  $\lambda_{\max}(A) \leq 1/\gamma$

# Lipschitz continuity

- $F$  is **Lipschitz continuous** with parameter  $L$  if it is single-valued on  $\text{dom } F$  and

$$\|F(x) - F(\hat{x})\|_2 \leq L\|x - \hat{x}\|_2 \quad \text{for all } x, \hat{x} \in \text{dom } F$$

- $F$  is **nonexpansive** if it is Lipschitz continuous with  $L = 1$

**Example:** any affine  $F(x) = Ax + b$  (parameter  $L = \|A\|_2$ )

## Relation to co-coercivity

- co-coercivity implies Lipschitz continuity (with  $L = 1/\gamma$ )
- Lipschitz continuity does not imply co-coercivity
- properties are equivalent for gradient of closed convex functions (page 1.15)



# Resolvent

the **resolvent** of an operator  $F$  is the operator

$$(I + \lambda F)^{-1} \quad (\text{with } \lambda > 0)$$

- inverse denotes the operator inverse:

$$y \in (I + \lambda F)^{-1}(x) \quad \iff \quad x - y \in \lambda F(y)$$

- graph of resolvent is a linear transformation of graph of  $F$ :

$$\text{gr}((I + \lambda F)^{-1}) = \begin{bmatrix} I & \lambda I \\ I & 0 \end{bmatrix} \text{gr}(F)$$

# Examples

**Subdifferential:** resolvent is proximal mapping

$$(I + \lambda \partial f)^{-1}(x) = \text{prox}_{\lambda f}(x)$$

follows from subgradient characterization of  $\text{prox}_{\lambda f}$  (page 4.7)

$$y = \text{prox}_{\lambda f}(x) \iff x - y \in \lambda \partial f(y)$$

**Monotone affine mapping:** resolvent of  $F(x) = Ax + b$  is

$$(I + \lambda F)^{-1}(x) = (I + \lambda A)^{-1}(x - \lambda b)$$

- inverse on right-hand side is standard matrix inverse
- $I + \lambda A$  is nonsingular for all  $\lambda \geq 0$  because  $A + A^T \geq 0$

# Monotonicity properties

- an operator is monotone if and only if its resolvent is firmly nonexpansive:

this follows from the matrix identity

$$\lambda \begin{bmatrix} 0 & I \\ I & 0 \end{bmatrix} = \begin{bmatrix} I & I \\ \lambda I & 0 \end{bmatrix} \begin{bmatrix} 0 & I \\ I & -2I \end{bmatrix} \begin{bmatrix} I & \lambda I \\ I & 0 \end{bmatrix}$$

and the expression of the graph of the resolvent on page [12.23](#)

- a monotone operator  $F$  is *maximal* monotone if and only

$$\text{dom}((I + \lambda F)^{-1}) = \mathbf{R}^n$$

follows from Minty's theorem on page [12.19](#)

# Outline

- primal–dual optimality conditions
- primal–dual hybrid gradient algorithm
- monotone operators
- **proximal point algorithm**

# Proximal point algorithm

**Monotone inclusion problem:** given maximal monotone  $F$ , find  $x$  such that

$$0 \in F(x)$$

this is equivalent to finding a fixed point of the resolvent  $R_t = (I + tF)^{-1}$  of  $F$ :

$$x = R_t(x) \iff x \in (I + tF)(x) \iff 0 \in F(x)$$

**Proximal point algorithm:** fixed point iteration

$$x_{k+1} = R_{t_k}(x_k)$$

**Proximal point algorithm with relaxation** (relaxation parameter  $\rho_k \in (0, 2)$ ):

$$x_{k+1} = x_k + \rho_k(R_{t_k}(x_k) - x_k)$$

# Convergence

if  $F^{-1}(0) \neq \emptyset$ , proximal point algorithm converges

- with constant  $t_k = t > 0$  and  $\rho_k = \rho \in (0, 2)$
- with  $t_k, \rho_k$  varying and bounded away from their limits, *i.e.*,

$$t_k \geq t_{\min} > 0, \quad 0 < \rho_{\min} \leq \rho_k \leq \rho_{\max} < 2 \quad \text{for all } k$$

proof relies on firm nonexpansiveness of resolvent

# Linear change of variables

make a change of variables  $x = Ay$ , with  $A$  nonsingular:

$$G(y) = A^T F(Ay)$$

- graph of  $G$  is

$$\text{gr}(G) = \begin{bmatrix} A^{-1} & 0 \\ 0 & A^T \end{bmatrix} \text{gr}(F)$$

- (maximal) monotonicity of  $G$  follows from (maximal) monotonicity of  $F$  and

$$\begin{bmatrix} A^{-1} & 0 \\ 0 & A^T \end{bmatrix}^T \begin{bmatrix} 0 & I \\ I & 0 \end{bmatrix} \begin{bmatrix} A^{-1} & 0 \\ 0 & A^T \end{bmatrix} = \begin{bmatrix} 0 & I \\ I & 0 \end{bmatrix}$$

## “Preconditioned” proximal point algorithm

$$y_{k+1} = (I + t_k G)^{-1}(y_k)$$

- $y_{k+1}$  is the solution  $y$  of the inclusion problem

$$\frac{1}{t_k}(y_k - y) \in A^T F(Ay)$$

- in the original coordinates  $x = Ay$ , this can be written as

$$\frac{1}{t_k}H(x_k - x) \in F(x)$$

where  $H = A^{-T}A^{-1}$  and  $x_k = Ay_k$

- we obtain a generalized proximal point update, with  $H > 0$  substituted for  $I$ :

$$x_{k+1} = (H + t_k F)^{-1}(Hx_k)$$

the purpose is often to make the resolvents cheaper, not preconditioning



# Proximal method of multipliers

the proximal point algorithm applied to

$$F(x, z) = \begin{bmatrix} 0 & A^T \\ -A & 0 \end{bmatrix} \begin{bmatrix} x \\ z \end{bmatrix} + \begin{bmatrix} \partial f(x) \\ \partial g^*(z) \end{bmatrix}$$

is known as the proximal method of multipliers

- basic iteration (without relaxation) is

$$(x_{k+1}, z_{k+1}) = (I + tF)^{-1}(x_k, z_k)$$

- $(x_{k+1}, z_{k+1})$  is the solution of the monotone inclusion with variables  $x, z$

$$0 \in \begin{bmatrix} 0 & A^T \\ -A & 0 \end{bmatrix} \begin{bmatrix} x \\ z \end{bmatrix} + \begin{bmatrix} \partial f(x) \\ \partial g^*(z) \end{bmatrix} + \frac{1}{t} \begin{bmatrix} x - x_k \\ z - z_k \end{bmatrix}$$

## Evaluation of the resolvent

- equivalent inclusion problem

$$0 \in \begin{bmatrix} 0 & 0 & A^T \\ 0 & 0 & -I \\ -A & I & 0 \end{bmatrix} \begin{bmatrix} x \\ y \\ z \end{bmatrix} + \begin{bmatrix} \partial f(x) \\ \partial g(y) \\ 0 \end{bmatrix} + \frac{1}{t} \begin{bmatrix} x - x_k \\ 0 \\ z - z_k \end{bmatrix}$$

- this is the optimality condition of the optimization problem (variables  $x, y$ )

$$\text{minimize } f(x) + g(y) + \frac{t}{2} \|Ax - y + (1/t)z_k\|_2^2 + \frac{1}{2t} \|x - x_k\|_2^2$$

(the augmented Lagrangian with an extra quadratic penalty term on  $x$ )

- from the minimizer  $(\hat{x}, \hat{y})$ , we make the update

$$x_{k+1} = \hat{x}, \quad z_{k+1} = z_k + t(A\hat{x} - \hat{y})$$

## PDHG and proximal point algorithm

apply “preconditioned” proximal point algorithm of page 12.29 with  $t_k = \tau$  and

$$H = \begin{bmatrix} I & -\tau A^T \\ -\tau A & (\tau/\sigma)I \end{bmatrix}$$

- $H$  is positive definite for  $\sigma\tau\|A\|_2^2 < 1$
- $x_{k+1}$  and  $z_{k+1}$  are the solution  $x, z$  of

$$\frac{1}{\tau} \begin{bmatrix} I & -\tau A^T \\ -\tau A & (\tau/\sigma)I \end{bmatrix} \begin{bmatrix} x_k - x \\ z_k - z \end{bmatrix} \in \begin{bmatrix} 0 & A^T \\ -A & 0 \end{bmatrix} \begin{bmatrix} x \\ z \end{bmatrix} + \begin{bmatrix} \partial f(x) \\ \partial g^*(z) \end{bmatrix}$$

- this simplifies to

$$0 \in \partial f(x) + \frac{1}{\tau}(x - x_k + \tau A^T z_k)$$

$$0 \in \partial g^*(z) + \frac{1}{\sigma}(z - z_k - \sigma A(2x - x_k))$$

can solve 1st inclusion for  $x$ ; substitute solution in 2nd inclusion and solve for  $z$

# PDHG and proximal point algorithm

$$0 \in \partial f(x) + \frac{1}{\tau}(x - x_k + \tau A^T z_k)$$

$$0 \in \partial g^*(z) + \frac{1}{\sigma}(z - z_k - \sigma A(2x - x_k))$$

- solution of the two inclusions is

$$x_{k+1} = (I + \tau \partial f)^{-1}(x_k - \tau A^T z_k)$$

$$z_{k+1} = (I + \sigma \partial g^*)^{-1}(z_k + \sigma A(2x_{k+1} - x_k))$$

- writing the solution in terms of prox operators gives the PDHG algorithm

$$x_{k+1} = \text{prox}_{\tau f}(x_k - \tau A^T z_k)$$

$$z_{k+1} = \text{prox}_{\sigma g^*}(z_k + \sigma A(2x_{k+1} - x_k))$$

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The convergence result on page 12.27 is Theorem 3 of this paper.