# Lecture 2 Piecewise-linear optimization

- piecewise-linear minimization
- $\ell_1$  and  $\ell_\infty$ -norm approximation
- examples
- modeling software

#### Linear and affine functions

**linear function:** a function  $f : \mathbf{R}^n \to \mathbf{R}$  is linear if

$$f(\alpha x + \beta y) = \alpha f(x) + \beta f(y) \qquad \forall x, y \in \mathbf{R}^n, \alpha, \beta \in \mathbf{R}$$

property: f is linear if and only if  $f(x) = a^T x$  for some a

**affine function:** a function  $f : \mathbf{R}^n \to \mathbf{R}$  is affine if

 $f(\alpha x + (1 - \alpha)y) = \alpha f(x) + (1 - \alpha)f(y) \qquad \forall x, y \in \mathbf{R}^n, \alpha \in \mathbf{R}$ 

property: f is affine if and only if  $f(x) = a^T x + b$  for some a, b

### **Piecewise-linear function**

 $f : \mathbf{R}^n \to \mathbf{R}$  is (convex) **piecewise-linear** if it can be expressed as

$$f(x) = \max_{i=1,...,m} (a_i^T x + b_i)$$

f is parameterized by m n-vectors  $a_i$  and m scalars  $b_i$ 



(the term *piecewise-affine* is more accurate but less common)

#### **Piecewise-linear minimization**

minimize 
$$f(x) = \max_{i=1,...,m} (a_i^T x + b_i)$$

• equivalent LP (with variables x and auxiliary scalar variable t)

minimize 
$$t$$
  
subject to  $a_i^T x + b_i \leq t, \quad i = 1, \dots, m$ 

to see equivalence, note that for fixed x the optimal t is t = f(x)

• LP in matrix notation: minimize  $\tilde{c}^T \tilde{x}$  subject to  $\tilde{A} \tilde{x} \leq \tilde{b}$  with

$$\tilde{x} = \begin{bmatrix} x \\ t \end{bmatrix}, \qquad \tilde{c} = \begin{bmatrix} 0 \\ 1 \end{bmatrix}, \qquad \tilde{A} = \begin{bmatrix} a_1^T & -1 \\ \vdots & \vdots \\ a_m^T & -1 \end{bmatrix}, \qquad \tilde{b} = \begin{bmatrix} -b_1 \\ \vdots \\ -b_m \end{bmatrix}$$

#### Minimizing a sum of piecewise-linear functions

minimize 
$$f(x) + g(x) = \max_{i=1,...,m} (a_i^T x + b_i) + \max_{i=1,...,p} (c_i^T x + d_i)$$

• cost function is piecewise-linear: maximum of mp affine functions

$$f(x) + g(x) = \max_{\substack{i=1,\dots,m\\j=1,\dots,p}} \left( (a_i + c_j)^T x + (b_i + d_j) \right)$$

• equivalent LP with m + p inequalities

minimize 
$$t_1 + t_2$$
  
subject to  $a_i^T x + b_i \le t_1, \quad i = 1, \dots, m$   
 $c_i^T x + d_i \le t_2, \quad i = 1, \dots, p$ 

note that for fixed x, optimal  $t_1$ ,  $t_2$  are  $t_1 = f(x)$ ,  $t_2 = g(x)$ 

• equivalent LP in matrix notation

$$\begin{array}{ll} \mbox{minimize} & \tilde{c}^T \tilde{x} \\ \mbox{subject to} & \tilde{A} \tilde{x} \leq \tilde{b} \end{array}$$

with

$$\tilde{x} = \begin{bmatrix} x \\ t_1 \\ t_2 \end{bmatrix}, \quad \tilde{c} = \begin{bmatrix} 0 \\ 1 \\ 1 \end{bmatrix}, \quad \tilde{A} = \begin{bmatrix} a_1^T & -1 & 0 \\ \vdots & \vdots & \vdots \\ a_m^T & -1 & 0 \\ c_1^T & 0 & -1 \\ \vdots & \vdots & \vdots \\ c_p^T & 0 & -1 \end{bmatrix}, \quad \tilde{b} = \begin{bmatrix} -b_1 \\ \vdots \\ -b_m \\ -d_1 \\ \vdots \\ -d_p \end{bmatrix}$$

## $\ell_{\infty}$ -Norm (Cheybshev) approximation

minimize  $||Ax - b||_{\infty}$ 

with  $A \in \mathbf{R}^{m \times n}$ ,  $b \in \mathbf{R}^m$ 

•  $\ell_{\infty}$ -norm (Chebyshev norm) of *m*-vector *y* is

$$||y||_{\infty} = \max_{i=1,\dots,m} |y_i| = \max_{i=1,\dots,m} \max\{y_i, -y_i\}$$

• equivalent LP (with variables x and auxiliary scalar variable t)

 $\begin{array}{ll} \text{minimize} & t\\ \text{subject to} & -t\mathbf{1} \leq Ax - b \leq t\mathbf{1} \end{array}$ 

(for fixed x, optimal t is  $t = ||Ax - b||_{\infty}$ )

• equivalent LP in matrix notation

minimize 
$$\begin{bmatrix} 0 \\ 1 \end{bmatrix}^T \begin{bmatrix} x \\ t \end{bmatrix}$$
  
subject to  $\begin{bmatrix} A & -\mathbf{1} \\ -A & -\mathbf{1} \end{bmatrix} \begin{bmatrix} x \\ t \end{bmatrix} \le \begin{bmatrix} b \\ -b \end{bmatrix}$ 

## $\ell_1$ -Norm approximation

minimize  $||Ax - b||_1$ 

•  $\ell_1$ -norm of *m*-vector *y* is

$$||y||_1 = \sum_{i=1}^m |y_i| = \sum_{i=1}^m \max\{y_i, -y_i\}$$

• equivalent LP (with variable x and auxiliary vector variable u)

minimize 
$$\sum_{i=1}^{m} u_i$$
  
subject to  $-u \leq Ax - b \leq u$ 

(for fixed x, optimal u is  $u_i = |(Ax - b)_i|$ , i = 1, ..., m)

• equivalent LP in matrix notation

minimize 
$$\begin{bmatrix} 0 \\ 1 \end{bmatrix}^T \begin{bmatrix} x \\ u \end{bmatrix}$$
  
subject to  $\begin{bmatrix} A & -I \\ -A & -I \end{bmatrix} \begin{bmatrix} x \\ u \end{bmatrix} \le \begin{bmatrix} b \\ -b \end{bmatrix}$ 

### **Comparison with least-squares solution**

histograms of residuals Ax - b, with randomly generated  $A \in \mathbf{R}^{200 \times 80}$ , for

$$x_{\rm ls} = \operatorname{argmin} \|Ax - b\|, \qquad x_{\ell_1} = \operatorname{argmin} \|Ax - b\|_1$$



 $\ell_1$ -norm distribution is wider with a high peak at zero

### **Robust curve fitting**

- fit affine function  $f(t) = \alpha + \beta t$  to m points  $(t_i, y_i)$
- an approximation problem  $Ax \approx b$  with

$$A = \begin{bmatrix} 1 & t_1 \\ \vdots & \vdots \\ 1 & t_m \end{bmatrix}, \qquad x = \begin{bmatrix} \alpha \\ \beta \end{bmatrix}, \qquad b = \begin{bmatrix} y_1 \\ \vdots \\ y_m \end{bmatrix}$$



- dashed: minimize ||Ax b||
- solid: minimize  $||Ax b||_1$

 $\ell_1\text{-norm}$  approximation is more robust against outliers

### Sparse signal recovery via $\ell_1$ -norm minimization

- $\hat{x} \in \mathbf{R}^n$  is unknown signal, known to be very sparse
- we make linear measurements  $y = A\hat{x}$  with  $A \in \mathbf{R}^{m \times n}$ , m < n

estimation by  $\ell_1$ -norm minimization: compute estimate by solving

minimize 
$$||x||_1$$
  
subject to  $Ax = y$ 

estimate is signal with smallest  $\ell_1$ -norm, consistent with measurements

equivalent LP (variables  $x, u \in \mathbf{R}^n$ )

$$\begin{array}{ll} \text{minimize} & \mathbf{1}^T u\\ \text{subject to} & -u \leq x \leq u\\ & Ax = y \end{array}$$

## Example

- exact signal  $\hat{x} \in \mathbf{R}^{1000}$
- 10 nonzero components



least-norm solutions (randomly generated  $A \in \mathbf{R}^{100 \times 1000}$ )



 $\ell_1$ -norm estimate is **exact** 

### **E**xact recovery

when are the following problems equivalent?

minimize	$\operatorname{card}(x)$	minimize	$\ x\ _1$
subject to	Ax = y	subject to	Ax = y

- card(x) is cardinality (number of nonzero components) of x
- depends on A and cardinality of sparsest solution of Ax = y

we say A allows **exact recovery** of k-sparse vectors if

$$\hat{x} = \underset{Ax=y}{\operatorname{argmin}} \|x\|_1$$
 when  $y = A\hat{x}$  and  $\operatorname{card}(\hat{x}) \le k$ 

- here,  $\operatorname{argmin} \|x\|_1$  denotes the unique minimizer
- a property of (the nullspace) of the 'measurement matrix' A

### 'Nullspace condition' for exact recovery

necessary and sufficient condition for exact recovery of k-sparse vectors<sup>1</sup>

$$|z^{(1)}| + \dots + |z^{(k)}| < \frac{1}{2} ||z||_1 \qquad \forall z \in \operatorname{nullspace}(A) \setminus \{0\}$$

here,  $z^{(i)}$  denotes component  $z_i$  in order of decreasing magnitude

$$|z^{(1)}| \ge |z^{(2)}| \ge \dots \ge |z^{(n)}|$$

- a bound on how 'concentrated' nonzero vectors in  $\operatorname{nullspace}(A)$  can be
- implies k < n/2
- difficult to verify for general A
- holds with high probability for certain distributions of random A

<sup>&</sup>lt;sup>1</sup>Feuer & Nemirovski (IEEE Trans. IT, 2003) and several other papers on compressed sensing.

### **Proof of nullspace condition**

#### notation

- x has support  $I \subseteq \{1, 2, \ldots, n\}$  if  $x_i = 0$  for  $i \notin I$
- |I| is number of elements in I
- $P_I$  is projection matrix on *n*-vectors with support *I*:  $P_I$  is diagonal with

$$(P_I)_{jj} = \begin{cases} 1 & j \in I \\ 0 & \text{otherwise} \end{cases}$$

• A satisfies the nullspace condition if

$$\|P_I z\|_1 < \frac{1}{2} \|z\|_1$$

for all nonzero z in  $\operatorname{nullspace}(A)$  and for all support sets I with  $|I| \leq k$ 

sufficiency: suppose A satisfies the nullspace condition

- let  $\hat{x}$  be k-sparse with support I (*i.e.*, with  $P_I \hat{x} = \hat{x}$ ); define  $y = A \hat{x}$
- consider any feasible x (*i.e.*, satisfying Ax = y), different from  $\hat{x}$
- define  $z = x \hat{x}$ ; this is a nonzero vector in  $\operatorname{nullspace}(A)$

$$\begin{aligned} |x||_{1} &= \|\hat{x} + z\|_{1} \\ &\geq \|\hat{x} + z - P_{I}z\|_{1} - \|P_{I}z\|_{1} \\ &= \sum_{k \in I} |\hat{x}_{k}| + \sum_{k \notin I} |z_{k}| - \|P_{I}z\|_{1} \\ &= \|\hat{x}\|_{1} + \|z\|_{1} - 2\|P_{I}z\|_{1} \\ &\geq \|\hat{x}\|_{1} \end{aligned}$$

(line 2 is the triangle inequality; the last line is the nullspace condition) therefore  $\hat{x} = \operatorname{argmin}_{Ax=y} \|x\|_1$  **necessity:** suppose A does not satisfy the nullspace condition

• for some nonzero  $z \in \operatorname{nullspace}(A)$  and support set I with  $|I| \leq k$ ,

$$\|P_I z\|_1 \ge \frac{1}{2} \|z\|_1$$

- define a k-sparse vector  $\hat{x} = -P_I z$  and  $y = A\hat{x}$
- the vector  $x = \hat{x} + z$  satisfies Ax = y and has  $\ell_1$ -norm

$$\begin{aligned} |x||_{1} &= \| -P_{I}z + z\|_{1} \\ &= \|z\|_{1} - \|P_{I}z\|_{1} \\ &\leq 2\|P_{I}z\|_{1} - \|P_{I}z\|_{1} \\ &= \|\hat{x}\|_{1} \end{aligned}$$

therefore  $\hat{x}$  is not the unique  $\ell_1$ -minimizer

### Linear classification

- given a set of points  $\{v_1, \ldots, v_N\}$  with binary labels  $s_i \in \{-1, 1\}$
- find hyperplane that strictly separates the two classes



homogeneous in a, b, hence equivalent to the linear inequalities (in a, b)

$$s_i(a^T v_i + b) \ge 1, \quad i = 1, \dots, N$$

### Approximate linear separation of non-separable sets



- penalty  $1 s_i(a_i^T v_i + b)$  for misclassifying point  $v_i$
- can be interpreted as a heuristic for minimizing #misclassified points
- $\bullet\,$  a piecewise-linear minimization problem with variables  $a,\,b$

equivalent LP (variables  $a \in \mathbf{R}^n$ ,  $b \in \mathbf{R}$ ,  $u \in \mathbf{R}^N$ )

$$\begin{array}{ll} \text{minimize} & \sum\limits_{i=1}^N u_i \\ \text{subject to} & 1-s_i(v_i^Ta+b) \leq u_i, \quad i=1,\ldots,N \\ & u_i \geq 0, \quad i=1,\ldots,N \end{array}$$

#### in matrix notation:

 $\begin{array}{ll} \text{minimize} & \begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix}^T \begin{bmatrix} a \\ b \\ u \end{bmatrix} \\ \text{subject to} & \begin{bmatrix} -s_1 v_1^T & -s_1 & -1 & 0 & \cdots & 0 \\ -s_2 v_2^T & -s_2 & 0 & -1 & \cdots & 0 \\ \vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\ -s_N v_N^T & -s_N & 0 & 0 & \cdots & -1 \\ 0 & 0 & -1 & 0 & \cdots & 0 \\ 0 & 0 & 0 & -1 & \cdots & 0 \\ \vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & 0 & \cdots & -1 \end{bmatrix} \begin{bmatrix} a \\ b \\ u_1 \\ u_2 \\ \vdots \\ u_N \end{bmatrix} \leq \begin{bmatrix} -1 \\ -1 \\ 0 \\ 0 \\ \vdots \\ 0 \end{bmatrix} \end{array}$ 

## Modeling software

modeling tools simplify the formulation of LPs (and other problems)

- accept optimization problem in standard notation (max,  $\|\cdot\|_1, \dots$ )
- recognize problems that can be converted to LPs
- express the problem in the input format required by a specific LP solver

#### examples of modeling packages

- AMPL, GAMS
- CVX, YALMIP (MATLAB)
- CVXPY, Pyomo, CVXOPT (Python)

## **CVX** example

minimize  $||Ax - b||_1$ subject to  $0 \le x_k \le 1$ ,  $k = 1, \dots, n$ 

#### MATLAB code

- between cvx\_begin and cvx\_end, x is a CVX variable
- after execution, x is MATLAB variable with optimal solution