EXTENSIONS OF SEMIDEFINITE PROGRAMMING METHODS FOR ATOMIC DECOMPOSITION

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ABSTRACT

We present an extension of recent semidefinite programming formulations for atomic decomposition over continuous dictionaries, with applications to continuous or ‘gridless’ compressed sensing. The dictionary considered in this paper is defined in terms of a general matrix pencil and is parameterized by a complex variable that varies over a segment of a line or circle in the complex plane. The main result of the paper is the formulation as a convex semidefinite optimization problem, and a simple constructive proof of the equivalence. The techniques are illustrated with a direction of arrival estimation problem, and an example of low-rank structured matrix decomposition.

Index Terms—Sparse signal reconstruction, low-rank matrix completion, compressed sensing, semidefinite optimization, array processing.

1. INTRODUCTION

Few optimization problems have attracted as much interest in recent years as the problem of minimizing the sum of a convex function and an $\ell_1$-norm regularization term. A general form of problems of this type is

$$\text{minimize} \quad f(\sum_{k=1}^r x_k a_k) + \sum_{k=1}^r |x_k|,$$

subject to $a_k \in D, \quad k = 1, \ldots, r,$

where $f$ is a convex function and $D$ is a set (or dictionary) of vectors in $\mathbb{C}^n$ or $\mathbb{R}^n$. The unknowns in problem (1) are the real or complex coefficients $x_k$, the vectors (or atoms) $a_1, \ldots, a_r$, selected from $D$, and the number $r$ of selected dictionary elements. If $D$ is a finite set, it can be represented by a matrix $D$ with the elements of $D$ as its columns, and the problem can be written as

$$\text{minimize} \quad f(Dx) + \|x\|_1.$$ 

This includes as special cases the Lasso problem [1], basis pursuit [2], noisy basis pursuit [3, 4], and numerous other applications [5–8].

When reviewing the literature on $\ell_1$-norm methods in signal processing [5, 9–11], it is striking that many of the underlying applications involve signals in continuous domains (time, space, or frequency domain), and the $\ell_1$-norm problems arise after discretizing and truncating an infinite dictionary. The discretization is used when no exact method for the continuous problem is known, or when the discretized problem is believed to be easier to solve numerically by convex optimization techniques.

It was recently noted that certain problems of the form (1) with infinite dictionaries can be exactly solved by semidefinite optimization. In particular, the authors of [12–17] consider $\ell_1$-norm minimization with dictionaries of vectors of undamped complex exponentials,

$$D = \{ \frac{1}{\sqrt{n}}(1, e^{j\omega}, e^{j2\omega}, \ldots, e^{j(n-1)\omega}) \mid \omega \in [0, 2\pi] \},$$

and use the fact that problem (1) is equivalent to the finite-dimensional convex optimization problem

$$\text{minimize} \quad f(y) + (\text{tr} X + z)/2$$

subject to $\begin{bmatrix} X & y \\ y^H & z \end{bmatrix} \succeq 0$

$X$ is Toeplitz.

The variables in this problem are $X \in \mathbb{H}^n$, $y \in \mathbb{C}^n$, $z \in \mathbb{R}$. (Here $\mathbb{H}^n$ denotes the set of Hermitian $n \times n$ matrices.)

A useful matrix extension of problem (1) results from allowing the variables $x_k$ to be vectors:

$$\text{minimize} \quad f(\sum_{k=1}^r a_k x_k^H) + \sum_{k=1}^r \|x_k\|_2$$

subject to $a_k \in D, \quad k = 1, \ldots, r.$

(3)

For example, if $D$ is the set of unit-norm vectors in $\mathbb{C}^n$, the problem can be shown to be equivalent to

$$\text{minimize} \quad f(Y) + \|Y\|_*$,$$

(4)

with variable $Y \in \mathbb{C}^{n \times m}$, where $\|Y\|_*$ denotes the trace norm or nuclear norm (sum of singular values). The equivalence can be seen by writing (4) in the equivalent form

$$\text{minimize} \quad f(\sum_{k=1}^r u_k v_k^H) + \sum_{k=1}^r (\|u_k\|_2^2 + \|v_k\|_2^2)/2$$

Research supported by NSF grants ECCS-1509789 and ECCS-1128817.
contains all vectors of the form $a_k = u_k/\sqrt{t_k}$, $x_k = \sqrt{t_k}v_k$, and eliminating the scalars $t_k$ by noting that $t_k + \|x_k\|^2/2$ is minimum for $t_k = \|x_k\|^2$.

The authors of [19, 20] consider (3) with the dictionary of complex exponentials (2) and show that the problem is equivalent to the convex problem

$$\min f(Y) + (\text{tr } X + \text{tr } Z)/2$$

subject to
$$X Y^H + Y X^H \geq 0$$
$$X$$ is Toeplitz,

with variables $X \in \mathbb{H}^n$, $Y \in \mathbb{C}^{n \times m}$, $Z \in \mathbb{H}^m$. This extension includes problem (1) as a special case with $m = 1$. Applications of (1) or (3) with the dictionary (2) arise frequently in signal processing, for example, in the estimation of line spectra, point source localization, and sensor array processing [5, 9, 10, 21].

In this paper we discuss extensions of the semidefinite programming formulations of (1) and (3) to a large class of dictionaries defined in terms of matrix pencils. In §2.1 we define the general dictionary and show that it includes as special cases the dictionary of complex exponentials (2) and subsets of the form

$$D = \left\{ \frac{1}{\sqrt{n}}(1, e^{j\omega}, e^{j2\omega}, \ldots, e^{j(n-1)\omega}) \mid |\omega - \alpha| \leq \beta \right\},$$

in which we restrict the frequencies to intervals $[\alpha - \beta, \alpha + \beta]$. In §2.2 we give a semidefinite programming formulation of (3) for the matrix pencil dictionary, and outline a simple and constructive proof of the equivalence. Section 3 illustrates the use of constrained dictionaries with two examples. Section 4 gives the concluding remarks.

2. SEMIDEFINITE OPTIMIZATION FORMULATION

2.1. Dictionary

Define two $(n-1) \times n$ matrices

$$F = \begin{bmatrix} 0 & I_{n-1} \end{bmatrix}, \quad G = \begin{bmatrix} I_{n-1} & 0 \end{bmatrix}$$

(7)

(where $I_{n-1}$ is the identity matrix of order $n-1$) and let $\mathcal{C}$ be the unit circle in the complex plane. Then the set

$$D = \{a \in \mathbb{C}^n \mid \|a\|_2 = 1, (\lambda G - F)a = 0, \lambda \in \mathcal{C}\}$$

(8)

contains all vectors of the form

$$a = \frac{c}{\sqrt{n}}(1, e^{j\omega}, e^{j2\omega}, \ldots, e^{j(n-1)\omega})$$

(9)

with $|c| = 1$. Up to a phase factor $c$, this is the dictionary of undamped complex exponentials in (2).

We generalize (8) in two ways. First, we allow $F$, $G$ to be arbitrary matrices, so $\lambda G - F$ is a general matrix pencil. Second, we define $\mathcal{C}$ to be a subset of the closed complex plane defined by a quadratic equality and inequality,

$$\mathcal{C} = \{\lambda \in \mathbb{C} \cup \{\infty\} \mid g_{\Phi}(\lambda, 1) = 0, g_{\Psi}(\lambda, 1) \leq 0\},$$

(10)

where $\Phi$ and $\Psi$ are $2 \times 2$ Hermitian matrices with $\det \Phi < 0$, and for a $2 \times 2$ Hermitian matrix $\Lambda$, we define

$$g_{\lambda}(\mu, \nu) = \begin{bmatrix} \mu & \nu \\ \nu^* & \mu^* \end{bmatrix} \Lambda \begin{bmatrix} \mu & \nu \\ \nu^* & \mu^* \end{bmatrix}.$$

When $\Phi_{11} = 0$ and $\Psi_{11} \leq 0$ we admit the point $\lambda = \infty$ in the set (10). (When $\lambda = \infty$, the condition $(\lambda G - F)a = 0$ in (8) is interpreted as $Ga = 0$.) The solution set of the equality $g_{\Phi}(\lambda, 1) = 0$ is a circle or straight line in the complex plane. For the purposes of this paper the most important example is

$$\Phi = \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix},$$

(11)

which parameterizes the unit circle. By adding the inequality in (10) we obtain segments of circles or lines [22]. For example, the dictionary of vectors (9) with $|\omega - \alpha| \leq \beta$ is defined by (11) and $\Psi_{11} = 0, \Psi_{12} = -c^{2\alpha}, \Psi_{22} = 2\cos \beta$.

2.2. Main result

The following theorem gives a convex formulation of problem (3), with the dictionary $D$ defined in (8) and (10). To simplify the handling of the point at infinity, we define

$$\mathcal{C} = \{\mu, \nu \neq 0 \mid g_{\Phi}(\mu, \nu) = 0, g_{\Psi}(\mu, \nu) \leq 0\}.$$

The pair $(\mu, \nu)$ is identified with a finite $\lambda = \mu/\nu$ if $\nu \neq 0$ and with $\lambda = \infty$ otherwise.

**Theorem 2.1** Let $F$, $G$ be matrices in $\mathbb{C}^{p \times n}$, and $\Phi$, $\Psi$ be Hermitian $2 \times 2$ matrices with $\det \Phi < 0$. Then the problem

$$\min f(Y) + (\text{tr } X + \text{tr } Z)/2$$

subject to
$$\mu_k G a_k = \nu_k F a_k, \quad k = 1, \ldots, r$$
$$\|a_k\|_2 = 1, \quad k = 1, \ldots, r$$

(12)

with variables $a_k \in \mathbb{C}^n$, $x_k \in \mathbb{C}^m$, $(\mu_k, \nu_k) \in \mathbb{C}^2$, and $r$, is equivalent to the semidefinite optimization problem

$$\min f(Y) + (\text{tr } X + \text{tr } Z)/2$$

s. t.
$$\Phi_{11} F X F^H + \Phi_{21} F X G^H + \Phi_{12} G X F^H + \Phi_{22} G X G^H = 0$$
$$\Psi_{11} F X F^H + \Psi_{21} F X G^H + \Psi_{12} G X F^H + \Psi_{22} G X G^H \leq 0$$

(13)

with variables $X \in \mathbb{H}^n$, $Y \in \mathbb{C}^{n \times m}$, and $Z \in \mathbb{H}^m$. 
Problem (5) is a special case with $F, G$ defined in (7), $\Phi$ defined in (11), and $\Psi = 0$. In this case, the matrix equality $FXFH = GXG^H$ states that the upper-left $(n-1) \times (n-1)$ block of $X$ equals its lower-right $(n-1) \times (n-1)$ block, or equivalently, $X$ is Toeplitz.

Our proof of theorem 2.1 relies on the following result, which is an immediate consequence of Corollary 1 in [23].

**Lemma 2.1** Suppose $X$ is positive semidefinite, has rank $r$, and satisfies the first two constraints in (13). Then $X$ can be factorized as $X = UU^H$ where $U \in \mathbb{C}^{n \times r}$ and

$$FU = W \text{diag}(\mu), \quad GU = W \text{diag}(\nu),$$

for some $W \in \mathbb{C}^{p \times r}$ and vectors $\mu, \nu \in \mathbb{C}^r$ that satisfy

$$(\mu_k, \nu_k) \in \mathbb{C}, \quad k = 1, \ldots, r.$$  

We omit the proof of this lemma, but emphasize that the proof outlined in [23] gives a simple constructive algorithm, based on singular value and Schur decompositions, for computing the factorization of $X$.

**Proof of Theorem 2.1.** We first show that (13) is a relaxation of (12). Suppose $a_1, \ldots, a_r, x_1, \ldots, x_r$ are feasible in (12). Without loss of generality assume that $x_k \neq 0$ for all $k$. Define $t_k = \sqrt{||x_k||_2}$ and

$$
\begin{bmatrix}
X & Y \\
Y^H & Z
\end{bmatrix} = \sum_{k=1}^r \begin{bmatrix}
t_k a_k \\
(1/t_k) x_k
\end{bmatrix} \begin{bmatrix}
t_k a_k \\
(1/t_k) x_k
\end{bmatrix}^H.
$$

Then $X, Y, Z$ are feasible in (13) with objective function equal to $f(\sum_k a_k x_k^H) + \sum_k ||x_k||_2$.

To establish the converse, we assume $X, Y, Z$ are feasible in (13) and show how to find a set of vectors $a_k, x_k$ that are feasible in (12) with objective function less than or equal to the objective of (13). Suppose $X$ has rank $r$. We can factorize $X$ as in lemma 2.1 and define $a_k = u_k/||u_k||_2$ and $\theta_k = ||u_k||_2$, where $u_k$ is the $k$th column of $U$. This gives a factorization

$$X = \sum_{k=1}^r \theta_k a_k a_k^H$$

where $\theta_k > 0$ and the vectors $a_k$ are linearly independent and satisfy the constraints in (12). The third constraint in (13) then implies that there exist vectors $x_1, \ldots, x_r$ such that

$$Y = \sum_{k=1}^r a_k x_k^H, \quad Z = \sum_{k=1}^r \frac{1}{\theta_k} x_k x_k^H.$$  

Substituting the expressions for $X, Y, Z$ in the objective gives

$$f(Y) + \frac{1}{2} (\text{tr}(X) + \text{tr}(Z)) \geq f(\sum_{k=1}^r a_k x_k^H) + \frac{1}{2} \sum_{k=1}^r (\theta_k + \frac{1}{\theta_k}) ||x_k||_2^2$$

$$\geq f(\sum_{k=1}^r a_k x_k^H) + \sum_{k=1}^r ||x_k||_2^2.$$  

The last line follows from the arithmetic-geometric mean inequality. \hfill \Box

### 3. EXAMPLES

#### 3.1. Direction of arrival estimation

The first example illustrates the interval constraints in dictionaries of the form (6). We consider a uniform linear array of $n = 50$ sensors. The signal arriving at the array is a superposition of a small number of planar waves arriving from different directions in $[-\pi/2, \pi/2]$. We take $2d/\lambda_c = 1$, where $d$ is the distance between the sensors and $\lambda_c$ the signal wavelength [21, §6.2]. When all sensor measurements are available, the directions of arrival can be estimated by classical methods, such as MUSIC and ESPRIT [21, 24, 25]. In this example, however, we assume that only a randomly selected subset of 30 sensors is used. Moreover the sensors are not omnidirectional. The 30 sensors are randomly partitioned in two groups of size 15. Sensors of group 1 measure signals arriving from directions in $[-\pi/2, \pi/2]$; sensors in group 2 measure signals arriving from directions $[-\pi/6, \pi/2]$. The output of the remaining 20 sensors is not used. To simplify notation we will assume the measurements are noise-free.

The direction of arrival (DOA) estimation problem can be put in the framework discussed in this paper by defining three dictionaries

$$D_j = \{(1, e^{i\omega_1}, e^{i\omega_2}, \ldots, e^{i(n-1)\omega}) \mid |\omega - \alpha_j| \leq \beta_j\},$$

for $j = 1, 2, 3$. Here $\omega = \pi \sin \theta$ is the spatial frequency associated with a direction of arrival $\theta$. The three intervals $[\alpha_j - \beta_j, \alpha_j + \beta_j]$ are the images of the intervals $[-\pi/2, -\pi/6], [-\pi/6, \pi/6], [\pi/6, \pi/2]$ after the transformation $\omega = \pi \sin \theta$. The DOA estimation problem can be formulated as an infinite-dimensional ‘basis-pursuit’ problem

$$\begin{align*}
\text{minimize} & \quad \sum_{j=1}^3 \sum_{k=1}^r |x_{jk}| \\
\text{subject to} & \quad y_j = \sum_{k=1}^r x_{jk} a_{jk}, \quad j = 1, 2, 3 \\
& \quad a_{jk} \in D_j, \quad k = 1, \ldots, r, \quad j = 1, 2, 3 \\
& \quad (y_1 + y_2) I_1 = b_1, \quad (y_2 + y_3) I_2 = b_2.
\end{align*}$$

Each of the variables $y_j$ is a linear combination of $r_j$ elements of the dictionary $D_j$, i.e., an $n$-vector that represents a sum of $r_j$ signals arriving from directions in the $j$th interval. In the last two constraints, $I_1 \subset \{1, 2, \ldots, n\}$ is an index set containing the indices of the sensors in group 1, and $I_2$ is an index set for group 2. The vector $b_1$ contains the 15 measurements at the sensors of group 1, and $b_2$ contains the measurements for group 2. The variables in the problem are the vectors $y_j$, $a_{jk}$, the coefficients $x_{jk}$, and the number of elements $r_j$ chosen from each dictionary.

Following the same idea as in theorem 2.1 we can cast the
problem as the SDP

\[
\begin{align*}
\text{min.} & \quad \sum_{j=1}^{3} (\text{tr}(X_j) + z_j)/2 \\
\text{s.t.} & \quad \begin{bmatrix} X_j & y_j \\ y_j^H & z_j \end{bmatrix} \succeq 0, \ j = 1, 2, 3 \\
& \quad X_1, X_2, X_3 \text{ are Toeplitz} \\
& \quad -e^{-j\omega_j} F X_j G^H - e^{j\omega_j} G X_j F^H \\
& \quad + 2 \cos \beta_j G X_j G^H \preceq 0, \ j = 1, 2, 3 \\
& \quad (y_1 + y_2)_1 = b_1, \ (y_2 + y_3)_2 = b_2
\end{align*}
\]

(16)

with variables \( X_j \in \mathbb{H}^n \), \( y_j \in \mathbb{C}^n \), \( z_j \in \mathbb{R} \), for \( j = 1, 2, 3 \). The matrices \( F, G \) are defined in (7).

Figure 1 shows an example. The red dots show the angles and magnitudes of 7 signals used to compute the measurement vectors \( b_1, b_2 \). The estimated angles and coefficients are shown with blue lines. The figure on the left-hand side shows that the angles and the coefficients are exactly recovered from the solution of (15) and (16). The right-hand plot shows the solution recovered from the SDP (16) if we omit the third set of constraints (which enforces the interval constraints). We also repeated the simulation with the same angles as in figure 1, different, randomly generated coefficients, and different random selections of the two sensor groups. The solution of the SDP (16) gave the exact answer in all instances, whereas the SDP without the interval constraints was successful in only about 27% of the instances.

3.2. Structured matrix decomposition

The second example is an application of theorem 2.1 with \( m > 1 \). We generate a 30 \( \times \) 30 matrix \( C = AB + N \) as a product of a 30 \( \times \) 3 matrix \( A \) with entries \( A_{ij} = \exp(j(i-1)\omega_j) \), for given values of \( \omega_1, \omega_2, \omega_3 \), and a randomly generated complex 3 \( \times \) 30 matrix \( B \) with entries from a normal distribution, plus a Gaussian noise matrix \( N \). The goal is to estimate the parameters \( \omega_j \) and the matrix \( B \) from the noisy measurements \( C \).

We compare two methods. In the first method we assume we are given a narrow interval that includes the parameters \( \omega_j \). In this case we define a dictionary

\[ D = \left\{ \frac{1}{\sqrt{n}} (1, e^{i\omega}, e^{2i\omega}, \ldots, e^{(n-1)i\omega}) \mid |\omega - \alpha| \leq \beta \right\} \]

\[ \text{min.} \quad \tau \left| \sum_{k=1}^{r} a_k x_k^H - C \right|_F + \sum_{k=1}^{r} ||x_k||_2 \\
\text{subject to} \quad a_k \in D, \ k = 1, \ldots, r. \]

(17)

with \( \tau \) a positive parameter. In the example, we use \( \alpha = 0, \beta = \pi/12 = 0.2618 \). Using theorem 2.1, the problem can be converted to the SDP

\[ \text{minimize} \quad \tau \|Y - C\|_F + (\text{tr}(X) + \text{tr}(Z))/2 \]

\[ \text{subject to} \quad \begin{bmatrix} X & Y \\ Y^H & Z \end{bmatrix} \succeq 0 \\
X \text{ is Toeplitz} \\
-FXG^H - GXF^H + 2 \cos \beta GXG^H \preceq 0. \]

In the second method, we omit the third constraint (imposing the interval), i.e., solve (17) with the dictionary (2). Figure 2 shows the two solutions, for independently tuned values of \( \tau \), with the estimates of \( \omega_i \) on the horizontal axis, and the norms of the vectors \( x_k \) on the vertical axis. As can be seen, adding the interval constraints allowed the method to identify the parameters \( \omega_i \), and thus \( \|x_k\| \), more accurately.

4. CONCLUSION

The proof technique and the factorization result in lemma 2.1 appear in the control literature on the generalized Kalman-Yakubovich-Popov lemma and its connections with SDP duality [23, 26–28]. In these applications, \( F \) and \( G \) are defined as

\[ F = \begin{bmatrix} A & B \end{bmatrix}, \ G = \begin{bmatrix} I & 0 \end{bmatrix}, \]

and the vectors \( a \) in (8) take the form \( a = ((\lambda I - A)^{-1}Bu, u) \). The link is important for two reasons. First, it suggests a range of similar applications in system and control theory, statistics, and numerical analysis, with different choices for the matrices \( F, G \), and the sets \( C \). Second, specialized techniques for solving SDPs derived from the Kalman-Yakubovich-Popov lemma, for example, by exploiting real symmetries and rank-one structure [29–32], will be useful in the development of fast solvers for SDPs of the form (13).
5. REFERENCES


